

# Leadership Statistics in Random Structures

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The largest component (“the leader”) in evolving random structures often exhibits universal statistical properties. This phenomenon is demonstrated analytically for two ubiquitous structures: random trees and random graphs. In both cases, lead changes are rare as the average number of lead changes increases quadratically with logarithm of the system size. As a function of time, the number of lead changes is self-similar. Additionally, the probability that no lead change ever occurs decays exponentially with the average number of lead changes.

E. Ben-Naim and P. L. Krapivsky, *Europhys. Lett.* **65**, 151 (2004).

## Problem

An ensemble of random growing trees (RT) or random growing graphs (RG), starting from  $N$  nodes.

## Questions

- What is the size of the leader, i.e., the largest component  $l(t, N)$ ?
- What is the number of lead changes  $L(t, N)$  as a function of time  $t$  and system size  $N$ ?
- What is the number of total lead changes  $L(N)$  as a function of system size?

## Motivation

- Data storage algorithms in computer science (RT) [1,2].
- Collision processes in gases (RT) [3].
- Random networks (RG) [4,5].
- Polymerization and Gelation (RG) [6].

# Random Trees

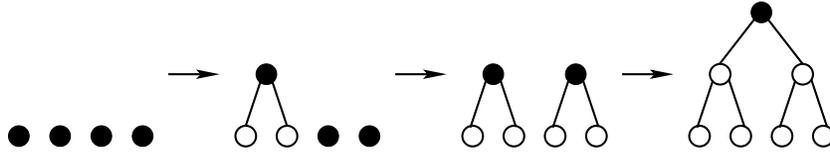


FIG. 1. Illustration of the tree merger process with system size  $N = 4$ .

- Start with  $N$  single-leaf trees.
- Pick two **trees** at random and merge them.

## Size Distribution

- Distribution of component of size  $k$  at time  $t$  is  $c_k(t)$
- Rate Equation

$$\frac{d}{dt}c_k = \sum_{i+j=k} c_i c_j - 2c c_k$$

- Solution subject to initial condition  $c_k(0) = \delta_{k,0}$ .

$$c_k(t) = \frac{t^{k-1}}{(1+t)^{k+1}}$$

- Exponential scaling distribution (asymptotically)

$$c_k(t) \simeq k_*^{-2} \Phi(k/k_*), \quad \Phi(x) = e^{-x}$$

- Typical size:  $k_* \sim t$ .

## Leadership Statistics

- **Leader Size**  $l(t, N)$ : Obtain leader size from cumulative distribution  $u_k = N \sum_{j=k}^{\infty} c_j$  using  $u_l = 1$ . Asymptotically,

$$l \simeq t \ln \frac{N}{t}.$$

- **Number of lead changes**  $L(t, N)$ : Obtain the rate by which the leader is surpassed from the rate of change in the cumulative distribution  $\frac{d}{dt}L = \frac{d}{dt}u_k|_{k=l}$ .

$$L(t, N) \simeq \ln t \ln N - \frac{1}{2}(\ln t)^2.$$

- **Unusual scaling form**: involves scaling function  $F(x) = x - \frac{1}{2}x^2$ .

$$L(t, N) = (\ln N)^2 F(x), \quad x = \frac{\ln t}{\ln N} = \frac{\ln k_*}{\ln N}$$

- **Total number of lead changes**  $L(N)$ : is obtained by noting that the condensation time, the time to form a single tree of size  $N$  is simply  $t_f \simeq N$ . Using  $x = 1$ ,

$$L(N) \simeq A(\ln N) \quad A = \frac{1}{2}.$$

# Numerical Simulations

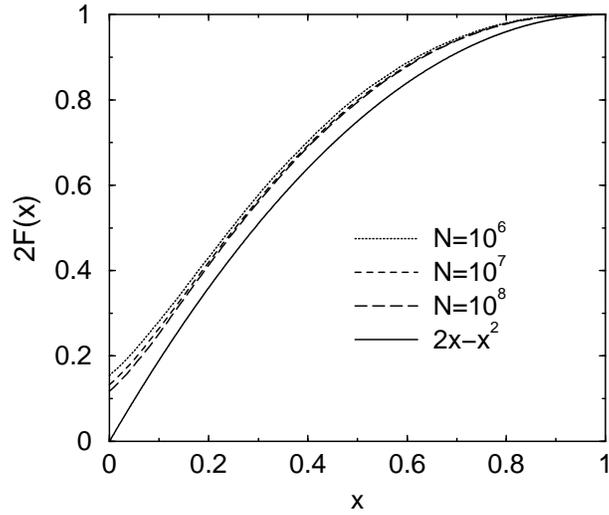


FIG. 2. The normalized time dependence of the number of lead changes for random trees,  $L(t, N)/L(N)$ , versus the scaling variable  $x = \ln t / \ln N$ . The simulation data, representing an average over  $10^3$  Monte Carlo runs, is compared with the theoretical prediction  $2F(x) = 2x - x^2$ .

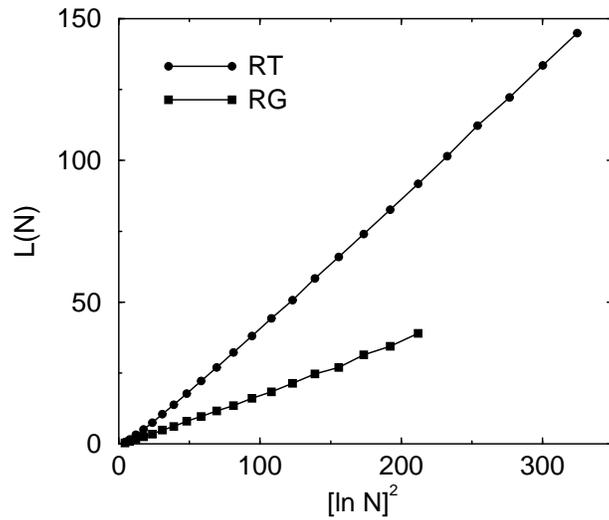


FIG. 3. The total number of lead changes  $L(N)$  versus the system size  $N$ . Shown are simulation results for Random Trees (RT) and Random Graphs (RG) representing an average over  $10^4$  realizations.

## Distribution of number of lead changes

- The probability  $P_n(t, N)$  that there are  $n$  lead changes at time  $t$  is Poissonian (assuming no correlations build up). Thus, it is characterized by the average number of lead changes  $L(t, n)$ .

$$P_n(t, N) = \frac{[L(t, N)]^n}{n!} e^{-L(t, N)}.$$

- The survival probability of the first leader  $S(t)$ , the probability that the initial leader is never overtaken, equals  $P_0(t, N)$

$$S(N) = \exp[-L] \simeq \exp[-A(\ln N)^2].$$

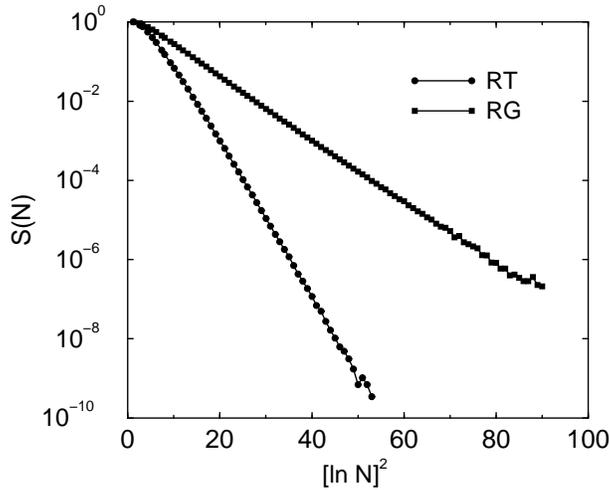


FIG. 4. The survival probability of the initial leader  $S(N)$  versus the system size  $N$ . The number of realizations was  $10^{10}$  and  $10^8$  for random trees and random graphs, respectively.

# Random Graphs

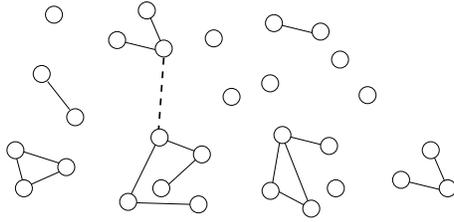


FIG. 5. Illustration of the random graph growth process.

- Start with  $N$  single-node graphs.
- Pick two **nodes** at random and merge their respective **graphs**.

## Size Distribution

- Distribution of components of size  $k$  at time  $t$ , is  $c_k(t)$ , satisfies the rate equation  $\frac{d}{dt}c_k = \frac{1}{2}\sum_{i+j=k} ij c_i c_j - k c_k$
- Solution subject to initial condition  $c_k(0) = \delta_{k,0}$

$$c_k(t) = \frac{(kt)^{k-1}}{k \cdot k!} e^{-kt}$$

- Scaling distribution (asymptotic)

$$c_k(t) \simeq k_*^{-5/2} \Phi(k/k_*), \quad \Phi(x) \propto x^{-5/2} e^{-x/2}.$$

- Gelation time:  $1 - t_g \sim N^{-1/3}$
- Typical size:  $k_* \sim (1 - t)^{-2}$  as  $t \rightarrow 1$ ; Giant comp. size  $\sim N^{2/3}$

# Leadership Statistics

- Obtained from size distribution as in the random tree case.

- **Leader size  $l(t, N)$ :**

$$l \simeq \frac{2}{(1-t)^2} \ln[N(1-t)^3].$$

- **Number of lead changes  $L(t, N)$ :**

$$L(t, N) \simeq 2 \ln N \ln \frac{1}{1-t} - 3 \left[ \ln \frac{1}{1-t} \right]^2.$$

- **Unusual scaling form:** involves  $F(x) = 2x - 3x^2$ .

$$L(t, N) \simeq (\ln N)^2 F(x), \quad x = \frac{\ln \frac{1}{1-t}}{\ln N} = \frac{1 \ln k_*}{2 \ln N}$$

- **Total number of lead changes  $L(N)$ :**

Obtained from  $1 - t_g \sim N^{-1/3}$  or  $x = 1/2$

$$L(N) \simeq A(\ln N) \quad A = \frac{1}{3}.$$

- **Probability no lead change occurs:**

$$S(N) \simeq \exp\left[-\frac{1}{3}(\ln N)^2\right]$$

## Conclusions

- Similar laws characterize random trees and random graphs.
- Lead changes are rare, their total number grows as  $(\ln N)^2$ .
- Unusual scaling behavior, with scaling variable  $\ln k_*/\ln N$ .
- Probability no lead change occur decays as  $\exp[-A(\ln N)^2]$ .

## References

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2. D. E. Knuth, *The Art of Computer Programming, vol. 3, Sorting and Searching* (Addison-Wesley, Reading, 1998).
3. R. van Zon, H. van Beijeren, and Ch. Dellago, *Phys. Rev. Lett.* **80**, 2035 (1998).
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